

# Human Reasoning, Computational Logic, and Ethical Decision Making

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# Inspiration

Luís Moniz Pereira and Ari Saptawijaya [2016]:  
Programming Machine Ethics

- ▶ Computational models of machine ethics
- ▶ Various ethical problems are implemented as logic programs
- ▶ Query for moral permissability
- ▶ However, the approach
  - ▶ does not provide a general method to account for ethical dilemmas
  - ▶ is not integrated into a cognitive theory about human reasoning

**We do not aim at suggesting a moral theory!**

The attempt of implementing a machine ethics, will help us understand human ethics and address the ambiguities that have not been sorted out so far. (Wallach and Allen, 2008)

# Trolley Problem (Foot [1967])

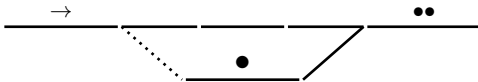
► Bystander Case



► Footbridge Case

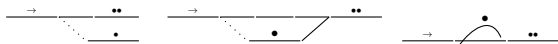


► Loop Case



Which action is morally permissible?

# Ethical Decision Principles in Trolley Problems



Bystander Case

Loop Case

Footbridge Case

	Bystander Case	Loop Case	Footbridge Case
Doctrine of double effect	change	-	-
Doctrine of triple effect	change	change	-
Maximize humans saved	change	change	throw down
action permissible say	85%	56%	12%

## Maximize the number of humans saved (Utilitarianism)

Could I save more humans by my action than humans that would be killed?

**Doctrine of double effect:** Killing is not permissible as *a means to save others*  
 If there were no human on the side track and I changed the switch then I would still save humans on the main track?

**Doctrine of triple effect:** *Intentional and direct* kill is not permissible  
 Could I avoid to intentionally and directly kill someone in order to save the others?

(Hauser, Cushman, Young, Kang-Xing Jin, Mikhail [2007]:  
 A Dissociation Between Moral Judgments and Justifications)

# Ethical Decision Making

Basic assumption

Humans construct models and reason with respect to them

An integrated computational cognitive theory must be able to consider

- ▶ actions with direct and indirect effects

- ▶ ethical principles

- ▶ conditional reasoning

*If I change the switch then I will save the humans on the main track*

- ▶ counterfactual or prefactual reasoning

Is a killing a side effect?

*If there were no human on the side track and I changed the switch then I would still save the humans on the main track*

This is ongoing work

# Towards an Integrated Computational Cognitive Theory

- ▶ Stenning, van Lambalgen [2009]  
Human Reasoning and Cognitive Science
- ▶ Hölldobler, Kencana Ramli [2009]  
Logic Programs under Three-Valued Łukasiewicz's Semantics

Normal logic programs  $\mathcal{P}$  are finite sets of

$$\begin{array}{ll} \text{Facts} & e \leftarrow \top \\ \text{Rules} & s \leftarrow e \wedge \neg ab_1 \qquad s \leftarrow t \wedge \neg ab_2 \\ \text{Assumptions} & ab_1 \leftarrow \perp \qquad ab_2 \leftarrow \perp \end{array}$$

Weak completion  $wc\mathcal{P}$  of program  $\mathcal{P}$

$$\{e \leftrightarrow \top, s \leftrightarrow (e \wedge \neg ab_1) \vee (t \wedge \neg ab_2), ab_1 \leftrightarrow \perp, ab_2 \leftrightarrow \perp\}$$

Least models under three-valued Łukasiewicz logic

$$\langle \{e, s\}, \{ab_1, ab_2\} \rangle$$

# Three-Valued Łukasiewicz Logic

truth values  $\{0, 1/2, 1\}$  (syntactically represented by  $\{\top, \perp, \bot\}$ )

negation  $\neg x \mapsto 1 - x$

(weak) disjunction  $x \vee y \mapsto \max(x, y)$

(weak) conjunction  $x \wedge y \mapsto \min(x, y)$

implication  $x \rightarrow y \mapsto \min(1, 1 - x + y)$

equivalence  $x \leftrightarrow y \mapsto 1 - |x - y|$

$\rightarrow$	0	1/2	1	$\leftrightarrow$	0	1/2	1
0	1	1	1	0	1	1/2	0
1/2	1/2	1	1	1/2	1/2	1	1/2
1	0	1/2	1	1	0	1/2	1

truth ordering  $0 <_t 1/2 <_t 1$

(total)

information ordering  $1/2 <_i 0$  and  $1/2 <_i 1$

(partial)

# Weak Completion Semantics of logic programs (WCS)

(Hölldobler and Kencana Ramli [2009])

**Semantic Operator**  $\Phi_{\mathcal{P}}(I) = \langle J^{\top}, J^{\perp} \rangle$  of ground program  $\mathcal{P}$ , where

$$\begin{aligned} J^{\top} &= \{A \mid A \leftarrow \text{Body} \in \mathcal{P} \text{ and } I(\text{Body}) = \top\} \\ J^{\perp} &= \{A \mid A \leftarrow \text{Body} \in \mathcal{P} \text{ and} \\ &\quad \text{for all } A \leftarrow \text{Body} \in \mathcal{P} \text{ we find } I(\text{Body}) = \perp\} \end{aligned}$$

**Least model** of weakly completed program  $\mathcal{P}$  = least fixed point of  $\Phi_{\mathcal{P}}$

$$\{e \leftarrow \top, s \leftarrow e \wedge \neg ab_1, s \leftarrow t \wedge \neg ab_2, ab_1 \leftarrow \perp, ab_2 \leftarrow \perp\}$$

	$\top$	$\perp$
$I$	$\emptyset$	$\emptyset$
$\Phi_{\mathcal{P}}(I)$	$\{e\}$	$\{ab_1, ab_2\}$
$\Phi_{\mathcal{P}}(\Phi_{\mathcal{P}}(I))$	$\{e, s\}$	$\{ab_1, ab_2\}$

$\Phi_{\mathcal{P}}(\Phi_{\mathcal{P}}(I))$  is a fixed point of  $\Phi_{\mathcal{P}}$



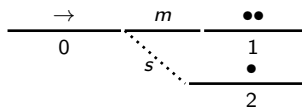
# Reasoning under Weak Completion Semantics

- ▶ Under WCS
  - ▶ represent a scenario as a logic program
  - ▶ compute the least model of the weak completion of the program
  - ▶ reason with respect to the least model
  - ▶ add skeptical abduction if necessary
  
- ▶ WCS is an integrated computational cognitive theory
  - ▶ suppression task
  - ▶ selection task
  - ▶ belief bias effect
  - ▶ syllogistic reasoning
  - ▶ spatial reasoning

How can we add **actions and causality** to WCS?

# Fluent Calculus (Hölldobler and Schneeberger [1990])

- ▶ **states** are represented as **multisets of fluents**
- ▶ states are changed by the **execution of actions**
- ▶ actions are specified by its **preconditions** and **direct effects**
- ▶ actions might have **indirect effects**, which can be computed by **ramifications**



$$\{t_0, c_0, m, h_1, h_1, h_2\} \xrightarrow{\text{change}} \{t_0, c_0, s, h_1, h_1, h_2\}$$

Fluents

$t_0 \quad t_1 \quad t_2 \quad m \quad s \quad c_0 \quad h_1 \quad h_2 \quad d$

Fluent terms

$t_0 \quad t_0 \circ c_0 \quad t_0 \circ c_0 \circ 1 \quad 1$  (unit)

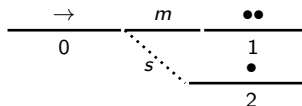
where  $\circ$  is an AC1-function symbol written infix

States

multisets of ethically irrelevant / relevant fluents

$(t_0 \circ c_0 \circ m \circ h_1 \circ h_1 \circ h_2, 1)$

# Actions



Agent

$action(1, 1, donothing, 1, 1) \leftarrow \top$

$action(m, 1, change, s, 1) \leftarrow \top$

Trolley

$action(t_0 \circ c_0 \circ m, 1, downhill, t_1 \circ c_0 \circ m, 1) \leftarrow \top$

$action(t_0 \circ c_0 \circ s, 1, downhill, t_2 \circ c_0 \circ s, 1) \leftarrow \top$

$action(t_1 \circ h_1, 1, kill, t_1, d) \leftarrow \top$

$action(t_2 \circ h_2, 1, kill, t_2, d) \leftarrow \top$

# Causality

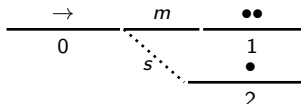
## Original fluent calculus

- ▶  $plan(X, P, Y)$  or  $causes(X, A, Y)$
- ▶ the execution of plan  $P$  transforms state  $X$  into state  $Y$ 
  - ▶ where a plan  $P$  is a sequence of actions
- ▶  $causes$  can be defined recursively on plans

## Problems:

- ▶ If a program  $\mathcal{P}$  contains recursive structures like lists or natural numbers then  $\Phi_{\mathcal{P}}$  is generally not continuous anymore  
Avoid recursive structures or restrict them to finite subsets
- ▶ There are infinitely many ground instances of  $causes(X, P, X)$ 
  - ▶ Consider as base case only finite scenarios
  - ▶ Consider only the states obtained by executing the actions of the agent
  - ▶ Compute successor states as ramifications wrt the actions of the trolley

# Weak Completion Semantics and Causality



## Base cases

$causes(\textit{do nothing}, t_0 \circ c_0 \circ m \circ h_1 \circ h_1 \circ h_2, 1) \leftarrow \top$

$causes(\textit{change}, t_0 \circ c_0 \circ s \circ h_1 \circ h_1 \circ h_2, 1) \leftarrow \top$

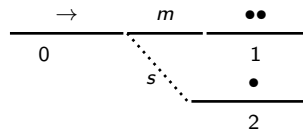
## Recursive case

$causes(A, E_1 \circ Z_1, E_2 \circ Z_2) \leftarrow \begin{aligned} &action(P_1, P_2, A', E_1, E_2) \\ &\wedge causes(A, P_1 \circ Z_1, P_2 \circ Z_2) \\ &\wedge \neg ab(A') \end{aligned}$

## Abnormalities

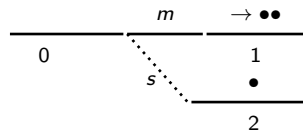
$ab(\textit{downhill}) \leftarrow \perp \quad ab(\textit{kill}) \leftarrow \perp$

# The Bystander Doing Nothing



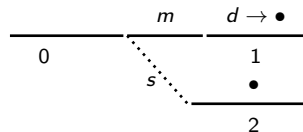
$causes(donothing, t_0 \circ c_0 \circ m \circ h_1 \circ h_1 \circ h_2, 1)$

$\Downarrow$  downhill



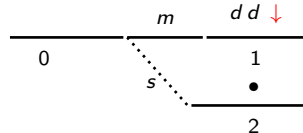
$causes(donothing, t_1 \circ c_0 \circ m \circ h_1 \circ h_1 \circ h_2, 1)$

$\Downarrow$  kill



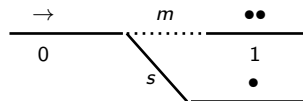
$causes(donothing, t_1 \circ c_0 \circ m \circ h_1 \circ h_2, d)$

$\Downarrow$  kill



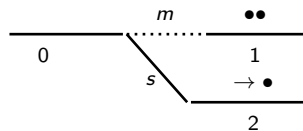
$causes(donothing, t_1 \circ c_0 \circ m \circ h_2, d \circ d)$

# The Bystander Changing the Switch



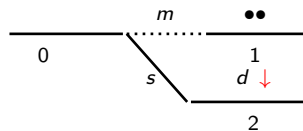
$causes(change, t_0 \circ c_0 \circ s \circ h_1 \circ h_1 \circ h_2, 1)$

↓ downhill



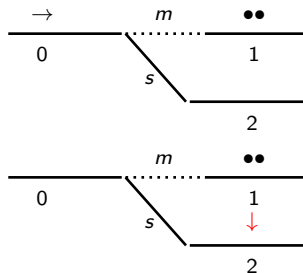
$causes(change, t_2 \circ c_0 \circ s \circ h_1 \circ h_1 \circ h_2, 1)$

↓ kill



$causes(change, t_2 \circ c_0 \circ s \circ h_1 \circ h_1, d)$

# The Bystander Changing Switch while Assuming Empty Side Track



$causes(change, t_0 \circ c_0 \circ s \circ h_1 \circ h_1 \circ c_2, 1)$

$\Downarrow$  downhill

$causes(change, t_2 \circ c_0 \circ s \circ h_1 \circ h_1 \circ c_2, 1)$



# Equational Theories

(Jaffar, Lassez, Maher [1984]:

A Theory of Complete Logic Programs with Equality)

$\mathcal{P}$  a (ground) normal logic program not containing the equality symbol

$\mathcal{E}$  a set of equations

$\equiv_{\mathcal{E}}$  finest congruence relation on the set of ground terms defined by  $\mathcal{E}$

$[t]$  congruence class defined by the ground term  $t$

Herbrand  $\mathcal{E}$ -universe quotient of the set of ground terms modulo  $\equiv_{\mathcal{E}}$

$[p(t_1, \dots, t_n)]$  abbreviation for  $p([t_1], \dots, [t_n])$

$[p(t_1, \dots, t_n)] = [q(s_1, \dots, s_m)]$  iff  $p = q$ ,  $n = m$ , and  $[t_i] = [s_i]$  for all  $i$

**Theorem** The weak completion of  $\mathcal{P}$  has a least  $\mathcal{E}$ -model under the three-valued Łukasiewicz logic

# Computing Least $\mathcal{E}$ -Models

Semantic Operator  $\Phi_{\mathcal{E}, \mathcal{P}}(I) = \langle J^\top, J^\perp \rangle$ , where

$$J^\top = \{[A] \mid A \leftarrow \text{Body} \in \mathcal{P} \text{ and } I(\text{Body}) = \top\}$$

$$J^\perp = \{[A] \mid A \leftarrow \text{Body} \in \mathcal{P} \text{ and for all } A' \text{ where} \\ A' \leftarrow \text{Body} \in \mathcal{P} \text{ with } [A] = [A'] \text{ we find } I(\text{Body}) = \perp\}$$

**Theorem**  $\Phi_{\mathcal{E}, \mathcal{P}}$  is monotonic.

It has a least fixed point. (by Knaster-Tarski Fixed Point Theorem)

Note that  $\Phi_{\mathcal{E}, \mathcal{P}}$  is not continuous in general.

$$q(1) \leftarrow \top \quad q(a \circ X) \leftarrow q(X) \quad r(1) \leftarrow \neg q(X)$$

Fixed point is reached after  $\omega + 1$  step, where  $\omega$  is the first limit ordinal.

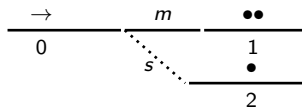
More results under the restriction to programs  $\mathcal{P}$  that are

- ▶ propositional,
- ▶ finite ground,
- ▶ or finite datalog programs with finite Herbrand  $\mathcal{E}$ -universe.

**Theorem**  $\Phi_{\mathcal{E}, \mathcal{P}}$  is continuous.

**Theorem** The least  $\mathcal{E}$ -model of the weak completion of  $\mathcal{P}$  is the least fixed point of  $\Phi_{\mathcal{E}, \mathcal{P}}$  and vice versa.

# Ethical Decision Making – The Bystander Case (1)



## Background Knowledge $\mathcal{P}_B$

$action(t_0 \circ c_0 \circ m, 1, downhill, t_1 \circ c_0 \circ m, 1) \leftarrow \top$   
 $action(t_0 \circ c_0 \circ s, 1, downhill, t_2 \circ c_0 \circ s, 1) \leftarrow \top$

$action(t_1 \circ h_1, 1, kill, t_1, d) \leftarrow \top$   
 $action(t_2 \circ h_2, 1, kill, t_2, d) \leftarrow \top$

$ab(downhill) \leftarrow \perp$   
 $ab(kill) \leftarrow \perp$

$causes(A, E_1 \circ Z_1, E_2 \circ Z_2) \leftarrow action(P_1, P_2, A', E_1, E_2)$   
 $\quad \wedge causes(A, P_1 \circ Z_1, P_2 \circ Z_2)$   
 $\quad \wedge \neg ab(A')$

## Ethical Decision Making – The Bystander Case (2)

*If I do nothing then the humans on the main track will be killed.* Yes

$\mathcal{P}_B$   
 $causes(donothing, t_0 \circ c_0 \circ m \circ h_1 \circ h_1 \circ h_2, 1) \leftarrow \top$

- ▶ Its least  $\mathcal{E}$ -model maps  $causes(donothing, t_1 \circ c_0 \circ m \circ h_2, d \circ d)$  to  $\top$

*If I change the switch then the humans on the main track will be saved.* Yes

*If I change the switch then the human on the side track will be killed.* Yes

$\mathcal{P}_B$   
 $causes(change, t_0 \circ c_0 \circ s \circ h_1 \circ h_1 \circ h_2, 1) \leftarrow \top$

- ▶ Its least  $\mathcal{E}$ -model maps  $causes(change, t_2 \circ c_0 \circ s \circ h_1 \circ h_1, d)$  to  $\top$

## Ethical Decision Making – The Bystander Case (3)

Changing the switch is preferable to do nothing as it will kill fewer humans. **Yes**

$\mathcal{P}_B$

$\text{causes}(\text{do nothing}, t_0 \circ c_0 \circ m \circ h_1 \circ h_1 \circ h_2, 1) \leftarrow \top$

$\text{causes}(\text{change}, t_0 \circ c_0 \circ s \circ h_1 \circ h_1 \circ h_2, 1) \leftarrow \top$

- ▶ Its least  $\mathcal{E}$ -model maps the following atoms to  $\top$

$\text{causes}(\text{do nothing}, t_1 \circ c_0 \circ m \circ h_2, d \circ d)$

$\text{causes}(\text{change}, t_2 \circ c_0 \circ s \circ h_1 \circ h_1, d)$

- ▶ Using

$\text{prefer}(A_1, A_2) \leftarrow \text{causes}(A_1, Z_1, D_1)$

$\wedge \text{causes}(A_2, Z_2, D_1 \circ d \circ D_2)$

$\wedge \neg \text{ab}_{\text{prefer}}(A_1)$

$\text{ab}_{\text{prefer}}(\text{change}) \leftarrow \perp$

$\text{ab}_{\text{prefer}}(\text{do nothing}) \leftarrow \perp$

In the least model the following atoms are mapped to  $\top$

$\text{causes}(\text{do nothing}, t_1 \circ c_0 \circ m \circ h_2, d \circ d)$

$\text{causes}(\text{change}, t_2 \circ c_0 \circ s \circ h_1 \circ h_1, d)$

The number of humans killed is minimized by changing the switch.

**Utilitarianism**

## Ethical Decision Making – The Bystander Case (4)

- ▶ *If there were no human on the side track and I changed the switch then I would still save the humans on the main track.* Yes

$\mathcal{P}_B$

$causes(change, t_0 \circ c_0 \circ s \circ h_1 \circ h_1 \circ c_2, 1) \leftarrow \top$

- ▶ Its least  $\mathcal{E}$ -model maps  $causes(change, t_2 \circ c_0 \circ s \circ h_1 \circ h_1 \circ c_2, 1)$  to  $\top$
- ▶ Using

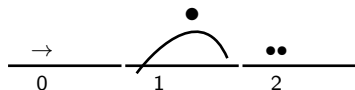
$permissible(change) \leftarrow prefer(change, donothing)$   
 $\quad \wedge causes(change, t_2 \circ c_0 \circ s \circ h_1 \circ h_1 \circ c_2, 1)$   
 $\quad \wedge \neg ab_{permissible}(change)$   
 $ab_{permissible}(change) \leftarrow \perp$

allows to conclude that changing the switch is permissible

### Doctrine of Double Effect

(Killing is permissible as a side effect but not as a means to save others)

# Ethical Decision Making – The Footbridge Case



## ► Base cases

$causes(donothing, t_0 \circ c_0 \circ c_1 \circ b_1 \circ h_2 \circ h_2, 1) \leftarrow \top$

$causes(throw, t_0 \circ c_0 \circ h_2 \circ h_2, d) \leftarrow \top$

## ► Is throwing the person from the bridge preferable to do nothing? **No**

$prefer(A_1, A_2) \leftarrow causes(A_1, Z_1, D_1)$

$\wedge causes(A_2, Z_2, D_1 \circ d \circ D_2)$

$\wedge \neg ab_{prefer}(A_1)$

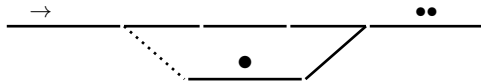
$ab_{prefer}(throw) \leftarrow intentional\_direct\_kill(throw)$

$intentional\_direct\_kill(throw) \leftarrow \top$

Pushing the person from the bridge is not permissible by

**Doctrine of Double Effect**

## Ethical Decision Making – The Loop Case



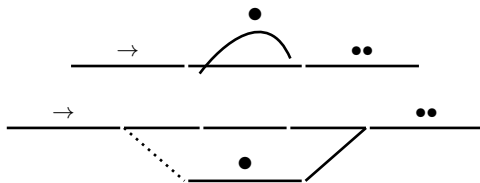
- ▶ *If I do nothing then the humans on the main track will be killed.* Yes
- ▶ *If I change the switch then the humans on the main track will be saved.* Yes  
*If I change the switch then the human on the side track will be killed.* Yes
- ▶ *If there were no human on the side track and I changed the switch then I would still save the humans on the main track.* No

Changing the switch is not permissible by

Doctrine of Double Effect



# Ethical Decision Making: Loop versus Footbridge Case



- ▶ Humans seem to distinguish the cases
- ▶ Throwing the person from the bridge is not permissible
- ▶ However, changing the switch is acceptable
- ▶ Direct versus indirect intentional kill

Could I avoid to intentionally and directly kill someone to save others?

**Doctrine of Triple Effect**

(Intentional and direct kill is not permissible.)

# Conclusion

- ▶ This is ongoing work
- ▶ We can solve all examples discussed in (Pereira, Saptawijaya 2017) uniformly in WCS with equality
- ▶ We are aiming at more general ethical rules
  - ▶ *If an action does something good and nothing abnormal is known then it is permissible.*
  - ▶ *A direct intentional kill is an abnormality.*
- ▶ Extension of WCS to more than three-valued Łukasiewicz logic